

Is converging bounded sequences  
really harder than  
converging binary sequences?

Claude Laflamme  
University of Calgary

## Historical Overview

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## Definition (Vojtáš (88))

$$\begin{aligned}
 \mathfrak{r} &= \min\{|\mathcal{A}| : \mathcal{A} \subseteq [\omega]^\omega \ \forall b \in 2^\omega \ \exists A \in \mathcal{A} \ \lim_{n \in A} b(n) \text{ exists.}\} \\
 &= \min\{|\mathcal{A}| : \forall B \ \exists A \in \mathcal{A} \ (A \subseteq^* B \text{ or } A \subseteq^* \omega \setminus B)\} \\
 &= \text{unsplitting number.}
 \end{aligned}$$

## Definition (Vojtáš (88))

$$\begin{aligned}
 \tau &= \min\{|\mathcal{A}| : \mathcal{A} \subseteq [\omega]^\omega \ \forall b \in 2^\omega \ \exists A \in \mathcal{A} \ \lim_{n \in A} b(n) \text{ exists.}\} \\
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$$\begin{aligned}
 \tau_\sigma &= \min\{|\mathcal{A}| : \mathcal{A} \subseteq [\omega]^\omega : \forall b \in \ell^\infty \ \exists A \in \mathcal{A} \ \lim_{n \in A} b(n) \text{ exists.}\} \\
 &= \min\{|\mathcal{A}| : \forall \langle B_n \rangle_n \ \exists A \in \mathcal{A} \ \forall n \ (A \subseteq^* B_n \text{ or } A \subseteq^* \omega \setminus B_n)\} \\
 &= \sigma - \text{unsplitting number.}
 \end{aligned}$$

## Proof of the Definition

Proof.

*Convergence*

$$\begin{array}{rcccccccc}
 & & & x_0 & x_1 & x_2 & x_3 & x_4 & x_5 & \cdots \\
 & & & \parallel & \parallel & \parallel & \parallel & \parallel & \parallel & \cdots \\
 \tau_\sigma & \chi_0 & = & 0 & 0 & 0 & 1 & 1 & 0 & \cdots \\
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**The verb “to reap” means “to split”, but as the letter  $\sigma$  has already been used, [2], we use  $\rho$ . (To be honest,  $\rho$  was suggested by Nyikos, who has a different reason for his choice of letter, and our term “to reap” is a back-formation. Furthermore, “to reap” does not really mean “to split”.)**

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# Early Questions

Question (Vojtáš 89)

$$Is \tau = \tau_\sigma?$$

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$$\text{Is } \mathfrak{r} = \mathfrak{r}_\sigma?$$

Question (Miller 82)

*Is  $cf(\mathfrak{r})$  uncountable?*

Observation

*$cf(\mathfrak{r}_\sigma)$  is uncountable.*

Remark (All easy...)

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$$\mathcal{A}_0 = \mathcal{A} \text{ and } \mathcal{A}_{n+1} = \{\pi_X(Y) : X, Y \in \mathcal{A}_n\}$$

So  $\bigcup \mathcal{A}_n$  has size  $\tau$ , and is “unsplitting below each member”.



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Now given  $\langle B_n \rangle_n$ , choose:

$$A_0 \in \mathcal{A}_0 : A_0 \subseteq^* B_0 \text{ or } A_0 \subseteq^* \omega \setminus B_0$$

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Now choose  $A \subseteq^* A_n$  for each  $n$  and this does unsplit  $\langle B_n \rangle_n$ .

## Theorem (Blass 93)

$$\max\{\mathfrak{d}, \mathfrak{r}\} \leq \mathfrak{hom}_n \leq \max\{\mathfrak{d}, \mathfrak{r}_\sigma\}.$$

## Definition

$$\mathfrak{d} = \min\{|\mathcal{D}| : \mathcal{D} \subseteq \omega^\omega \ \forall g \in \omega^\omega; \exists f \in \mathcal{D} \ f \geq^* g.\}$$

$$\mathfrak{hom}_n = \min\{|\mathcal{A}| : \mathcal{A} \subseteq [\omega]^\omega : \forall h : [\omega]^n \rightarrow 2 \exists A \in \mathcal{A} \ h \upharpoonright [A]^n =^* \text{cte.}\}$$

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Proof  $\mathfrak{hom}_2 \leq \max\{\mathfrak{d}, \mathfrak{r}_\sigma\}$ .

For  $h$  in a dominating family  $\mathcal{D}$ ,  $X$  in a  $\sigma$ -unsplitting family  $\mathcal{R}$  and  $Y \in \pi_X(\mathcal{R})$ , choose

$$H(h, X, Y) \subseteq Y \text{ infinite so that } x < y \implies h(x) < y$$

. Then  $\{H(h, X, Y) : h \in \mathcal{D}, X \in \mathcal{R}, Y \in \pi_X(\mathcal{R})\}$  works for  $\mathfrak{hom}_2$ . □

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If  $h \upharpoonright [A]^2 = cte$ , then  $\chi_n \upharpoonright A =^* cte$  for all  $n$ . □

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$$C_{\alpha, \gamma}^0 = \omega$$

$$C_{\alpha, \gamma}^{n+1} = \begin{cases} A_{\pi_\alpha(g_\gamma(n))} & \text{if this set is almost contained in } C_{\alpha, \gamma}^n, \\ C_{\alpha, \gamma}^n & \text{otherwise.} \end{cases}$$

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In the end, let  $C_{\alpha, \gamma}$  be an infinite pseudointersection of the  $C_{\alpha, \gamma}^n$ .

Show that the sets  $C_{\alpha, \gamma}$  form a  $\sigma$ -unsplitting family: given  $\langle D_n \rangle_n$ :

$$E = \{F \subset \tau : \forall n \forall \beta \in F \exists \delta \in F A_\delta \subseteq^* A_\beta \cap D_n \text{ or } A_\delta \subseteq^* A_\beta \setminus D_n\}$$





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- $\delta < \tau_\sigma$  in random real model.
- $\tau < u$  in Goldstern-Shelah model.
- Finite support iteration forces  $non(\mathcal{M}) \leq \tau$ , so cannot yield  $\tau = \aleph_1 < \tau_\sigma = \aleph_2$ .



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## Conjecture

*Andrzej will show uncountable support iteration won't work either....*

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- So the ultrafilter is a  $P$ -point ( $\mathfrak{u} = \tau < \mathfrak{d}$ ), thus  $\tau = \tau_\sigma$ .



$$\begin{array}{c}
 \mathfrak{r}_\sigma = \text{hom}_2 \\
 \begin{array}{c}
 \text{u} \quad | \\
 \quad \quad | \\
 \quad \quad | \\
 \quad \quad \mathfrak{r} \\
 \quad \quad | \\
 \quad \quad \mathfrak{d} \\
 \quad \quad | \\
 \quad \quad \mathfrak{b}
 \end{array} \\
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 | \\
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# $\mathfrak{fr}$ and $\mathfrak{fr}_\sigma$

## Definition (Brendle 98)

$\mathfrak{fr} := \min\{|\mathcal{A}| : \mathcal{A} \text{ consists of partitions of } \omega \text{ into finite sets, and no single } X \subseteq \omega \text{ splits every element of } \mathcal{A}\}$

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## Theorem (Brendle 98)

$$\mathfrak{r} = \min\{\mathfrak{d}, \mathfrak{r}\} \text{ and } \min\{\mathfrak{d}, \mathfrak{r}_\sigma\} = \mathfrak{r}_\sigma.$$

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## Theorem (Aubrey 04)

$\min\{\mathfrak{d}, \mathfrak{r}\} = \min\{\mathfrak{d}, \mathfrak{r}_\sigma\}$  and thus :

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## Question

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